

Addition by Fourier transform

Carsten Urbach

This corresponds to problem 5.6 in Nielsen & Chuang. The original paper is (Draper 2000). Which quantum circuit can be used to perform the computation

$$|x\rangle \rightarrow |x + y \pmod{2^n}\rangle$$

with $0 \leq x < 2^n$ and a constant **integer** y .

We exploit the general idea

$$x + y = \log(e^x e^y)$$

where the exponentiation is de facto performed by a Fourier trafo and the logarithm by the inverse trafo.

Fourier transforming the state $|x\rangle$ with n bits, leads to the following product representation

$$|x\rangle = |x_n x_{n-1} \dots x_1\rangle \rightarrow \frac{1}{2^n} (|0\rangle + e^{2\pi i 0 \cdot x_1} |1\rangle) (|0\rangle + e^{2\pi i 0 \cdot x_2 x_1} |1\rangle) \dots (|0\rangle + e^{2\pi i 0 \cdot x_n \dots x_1} |1\rangle)$$

where we use the notation

$$x = x_1 2^0 + x_2 2^1 + \dots + x_n 2^{n-1}$$

and

$$0.x_l \dots x_1 \equiv \frac{x_l}{2} + \frac{x_{l-1}}{2^2} + \dots + \frac{x_1}{2^l}.$$

Now, we apply a phase shift $R_\theta(\theta)$ to each qubit

$$R_z \equiv \begin{pmatrix} 1 & 0 \\ 0 & \exp(i\theta) \end{pmatrix}.$$

We apply R_θ with $\theta_j = 2\pi y / 2^{n-(j-1)}$ to qubit j where $1 \leq j \leq n$. For y we can also write

$$y = y_1 2^0 + y_2 2^1 + \dots + y_n 2^{n-1}.$$

Thus,

$$\exp(2\pi i y / 2^{n-j+1}) = \prod_{k=0}^{n-1} \exp(2\pi i y_{k+1} 2^{j-1-n+k}).$$

Since $\exp(2\pi i y_k l) = 1$ for positive integer l , this reduces to (recall $y_k \in \{0, 1\}$)

$$\exp(2\pi i y / 2^{n-j+1}) = \prod_{k=0}^{n-j} \exp(2\pi i y_{k+1} 2^{j-1-n+k}).$$

The n th qubit gets multiplied with $\exp(i\theta_n)$ with $\theta_n = 2\pi y / 2^1$. Thus, we need to compute

$$\exp(2\pi i x_1 / 2) \cdot \exp(2\pi i y_1 / 2) = \exp(2\pi i (x_1 + y_1) / 2).$$

Similarly, for the j th qubit one gets

$$\exp(2\pi i (x_1 / 2^{n-j+1} + x_2 / 2^{n-j} + \dots)) \cdot \exp(2\pi i (y_1 / 2^{n-j+1} + y_2 / 2^{n-j} + \dots)) = \exp(2\pi i ((x_1 + y_1) / 2^{n-j+1} + (x_2 + y_2) / 2^{n-j} + \dots))$$

which implements the addition mod n operation in this binary fraction.

Now apply the inverse Fourier trafo and it is easy to see that this transforms back to the state $|x + y \bmod n\rangle$.

For the practical implementation we first need the phase shift operators, which is up to a phase identical to R_z :

```
Rtheta <- function(bit, theta=0.) {  
  return(methods::new("sqgate", bit=as.integer(bit),  
    M=array(as.complex(c(1, 0, 0, exp(1i*theta))),  
      dim=c(2,2)), type="Rt"))  
}
```

With this one can write the desired function on state x .

```
addbyqft <- function(x, y) {  
  n <- x@nbits  
  z <- qsimulatR::qft(x)  
  for(j in c(1:n)) {  
    z <- Rtheta(bit=j, theta = 2*pi*y/2^(n-j+1)) * z  
  }  
  z <- qft(z, inverse=TRUE)  
  return(invisible(z))  
}
```

Examples

```
x <- qstate(5, basis=as.character(seq(0, 2^5-1)))  
x
```

```
( 1 ) * 0
```

```
z <- addbyqft(x, 3)  
z
```

```
( 1 ) * 3
```

```
z <- addbyqft(z, 5)  
z
```

```
( 1 ) * 8
```

```
z <- addbyqft(z, 30)  
z
```

```
( 1 ) * 6
```

References

Draper, Thomas G. 2000. "Addition on a Quantum Computer." *arXiv Preprint Quant-Ph/0008033*.